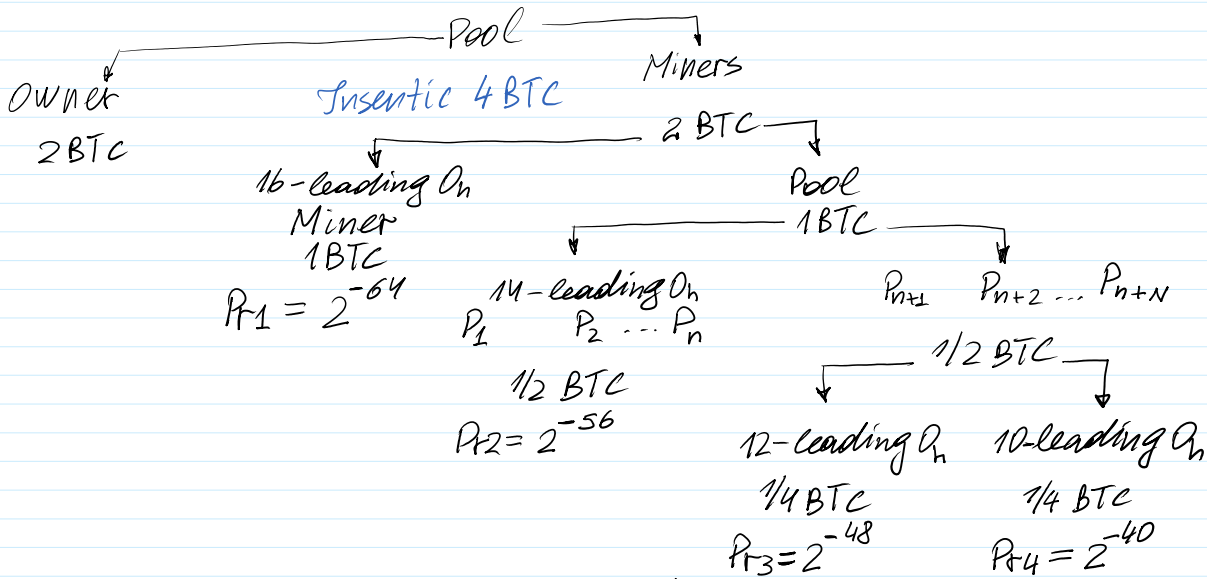


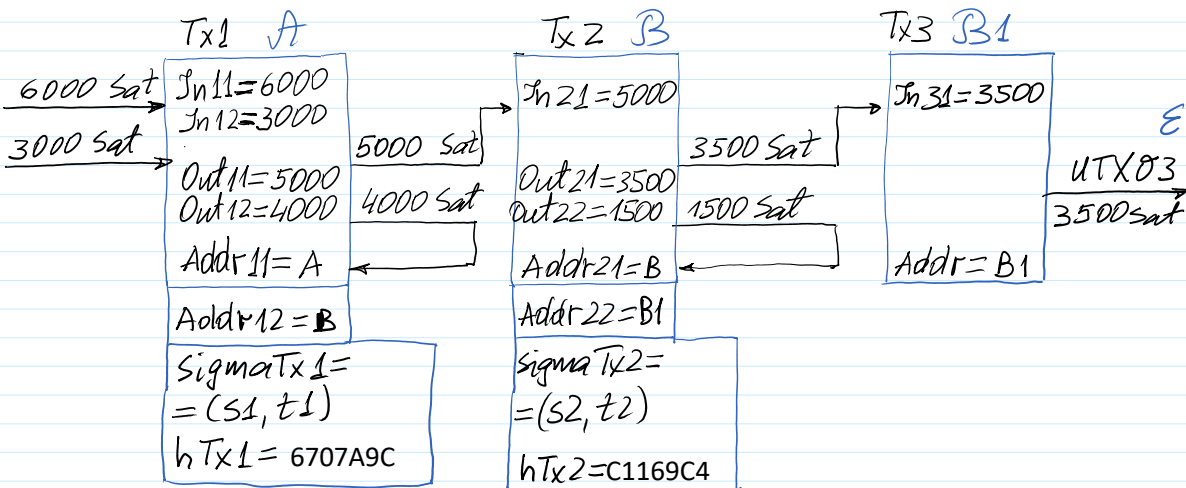
Shares

Pr { to compute h-value with 16 leading hex n. } = $\frac{2^{192}}{2^{256}} =$
 $= 2^{192} \cdot 2^{-256} = 2^{192-256} = 2^{-64}$ // If $Pr = \frac{1}{2} = 2^{-1} \Rightarrow$ 2 trials in average.



Assume student has $1 \text{ Th/s} = 2^{40} \text{ h/s}$.
 Let the number of competitors is 2^{20} .
 Then the reward for mined 10-leading Q_n
 $1 \cdot 2^{-2} / 2^{20} = 2^{-22} \text{ BTC}$.

Block mining



$Tx1 = '1 : \text{In} = 9000 \parallel \text{Out}_1 = 5000 \parallel \text{Out}_2 = 4000 \parallel B \parallel A'$
 $Tx2 = '2 : \text{In} = 5000 \parallel \text{Out}_1 = 3500 \parallel \text{Out}_2 = 1500 \parallel B1 \parallel B'$

$Tx3 = '3 : In=3500 || Out1=3500 || Out2=0 || E || B1'$

3 Transactions: 1 Input --> 2 Output

Merkle Tree 3 --> 1

$\gg h00 = h28(Tx1) \quad \% h00 = 'ansv'$

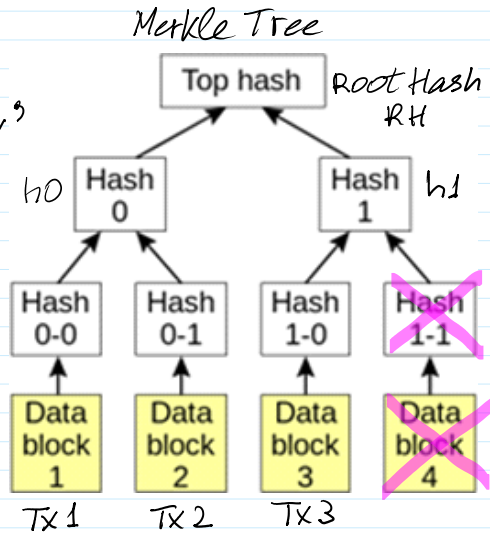
$\gg h01 = h28(Tx2)$

$\gg h10 = h28(Tx3)$

$\gg h0 = h28(h00 || h01)$

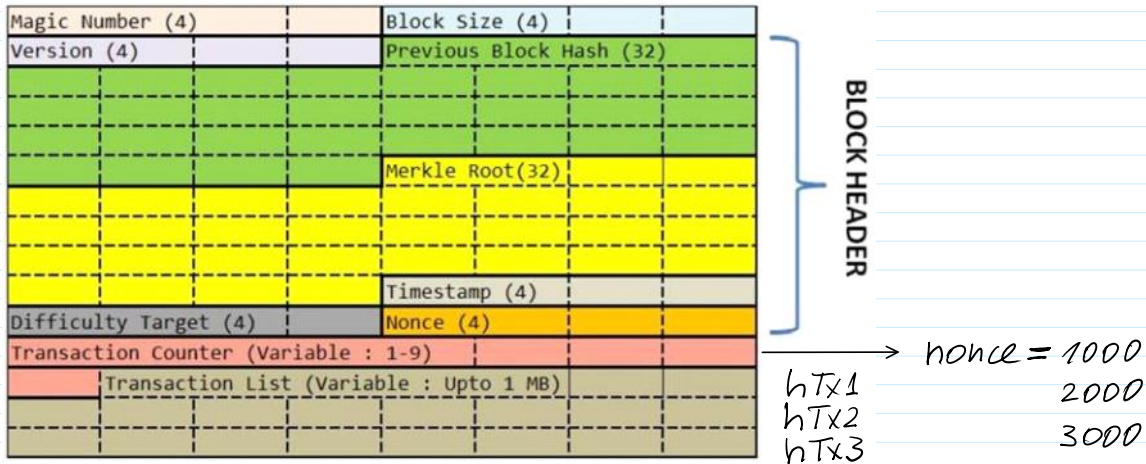
$\gg h1 = h10$

$\gg RH = h28(h0 || h1)$



$\gg hTx1 = \text{dec2hex}(h1Sig)$
 $hTx1 = 6707A9C$

$\gg hTx2 = \text{dec2hex}(h2Sig)$
 $hTx2 = C1169C4$



$\gg PrBlh = h28('nonce=1020')$
 $PrBlh = OCAF06F$

Difficulty target $DT = \phi H H H H H H H$

Block 1: $PrBlh: PrBlh=OCAF06F$

$Rh: Rh=89B5B94$

$hTx1=6707A9C$

$hTx2=C1169C4$

nonce : nonce=1000
 nonce=2000

Person No 1 in AIS list
 No2

nonce :	nonce=1000	Person No 1 in AIS list
	nonce=2000	No2
	nonce=3000	No3
	

Block numbers corresponds to your numbers in AIS list.

'Bl1 : PrBl = ... || Rh1 = ... || hTx1 = ... || hTx2 = ... || nonce = 1000'

```
>> Bl1M=h28('Bl1:PrBlh=0CAF06F| |Rh=89B5B94| |hTx1=6707A9C| |hTx2=C1169C4| |nonce=1000')
Bl1M = D99C1E7
```

```
>> Bl1M=h28('Bl1:PrBlh=0CAF06F| |Rh=89B5B94| |hTx1=6707A9C| |hTx2=C1169C4| |nonce=2018')
Bl1M = 06B0772
```

h28 : computes h-value of 28 bits

According to the difficulty target the adequate mined block must have leading hexadecimal 0 or 4 leading 0_b bits.

All available h-values with 28 bits length is equal to

```
>> 2^28 ans = 268435456
```

The number of adequate mined values is equal to any
28 - 4 = 24 bits values

```
>> 2^24 ans = 16777216
```

The probability to mine a block is the following:

$$Pr = \frac{\text{adequat number of values}}{\text{total number of values}} = \frac{2^{24}}{2^{28}} = 2^{-4} = \frac{1}{16}$$

<http://crypto.fmf.ktu.lt/xdownload/>

- [octave-6.3.0-w64-installer.exe](#)

Updated functions: h24(), hd24(), AES128().

- [Octave Stud 2021.10 Updated.7z](#)